

HDP

(19) Gaussian Width and Random Projection

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前情回顾

Analyze $\mathbb{E} \sup_{t \in T} X_t$ for Gaussian Process

1. Slepian's inequality: Small fluctuation leads to small expectation.
2. Sudakov-Fernique's inequality: remove variance requirement (only expectation results)
3. Gordon's inequality: two-dim extension
4. *Sudakov's minoration inequality: lower bound, Geometric

Random Process and Operator Norm

$$\|A\| = \max_{u \in S^{n-1}, v \in S^{m-1}} \langle Au, v \rangle = \max_{(u,v) \in T} X_{uv}$$

Proof based on covering number

$$\mathbb{E} \|A\| \leq \sqrt{m} + C\sqrt{n}$$

Proof based on Gaussian process

$$\mathbb{E} \|A\| \leq \sqrt{m} + \sqrt{n}.$$

Intuition: Apply Sudakov-Fernique's inequality

Sudakov-Fernique's inequality leads to Sudakov's minoration inequality.

Therefore the bound is tighter than covering number – based bound.

(just an intuition, minoration bound is in fact lower bound)

Gaussian width

Definition 7.5.1. The *Gaussian width* of a subset $T \subset \mathbb{R}^n$ is defined as

$$w(T) := \mathbb{E} \sup_{x \in T} \langle g, x \rangle \quad \text{where} \quad g \sim N(0, I_n).$$

a special width which is invariant under affine and convex operator

Note: g (symmetric) has an expectation.

Therefore, the asymmetric T suffers a small $w(T)$.

$$\frac{1}{\sqrt{2\pi}} \cdot \text{diam}(T) \leq w(T) \leq \frac{\sqrt{n}}{2} \cdot \text{diam}(T).$$

We can also replace g to other random vectors, e.g., spherical.

Gaussian width

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Some examples:

$$\ell_2 \text{ ball} \quad w(S^{n-1}) = w(B_2^n) = \mathbb{E} \|g\|_2 = \sqrt{n} \pm C,$$

$$\ell_\infty \text{ ball} \quad w(B_\infty^n) = \mathbb{E} |g_1| \cdot n = \sqrt{\frac{2}{\pi}} \cdot n.$$

$$\ell_1 \text{ ball} \quad c\sqrt{\log n} \leq w(B_1^n) \leq C\sqrt{\log n}.$$

$$\ell_p \text{ ball} \quad w(B_p^n) \leq C\sqrt{p'} n^{1/p'}.$$

$$\text{Finite set} \quad w(T) \leq C\sqrt{\log |T|} \cdot \text{diam}(T).$$

Asymmetric set suffers.

Stable dimension

$$w(T) := \mathbb{E} \sup_{x \in T} \langle g, x \rangle \quad \text{where} \quad g \sim N(0, I_n).$$

$$h(T)^2 := \mathbb{E} \sup_{t \in T} \langle g, t \rangle^2, \quad \text{where} \quad g \sim N(0, I_n).$$

Definition 7.6.2 (Stable dimension). For a bounded set $T \subset \mathbb{R}^n$, the *stable dimension* of T is defined as

$$d(T) := \frac{h(T - T)^2}{\text{diam}(T)^2} \asymp \frac{w(T)^2}{\text{diam}(T)^2}.$$

$d(T) \leq n$; 球上二者相等; 有限集合 $\log |T|$

Definition 7.6.7 (Stable rank). The *stable rank* of an $m \times n$ matrix A is defined as

$$r(A) := \frac{\|A\|_F^2}{\|A\|^2}.$$

Gaussian complexity

$$w(T) := \mathbb{E} \sup_{x \in T} \langle g, x \rangle \quad \text{where} \quad g \sim N(0, I_n).$$

Definition 7.6.8. The *Gaussian complexity* of a subset $T \subset \mathbb{R}^n$ is defined as

$$\gamma(T) := \mathbb{E} \sup_{x \in T} | \langle g, x \rangle | \quad \text{where} \quad g \sim N(0, I_n).$$

Exercise 7.6.9 (Gaussian width vs. Gaussian complexity). ☕☕☕ Consider a set $T \subset \mathbb{R}^n$ and a point $y \in T$. Show that

$$\frac{1}{3} [w(T) + \|y\|_2] \leq \gamma(T) \leq 2 [w(T) + \|y\|_2]$$

Random Projection

Consider the random projection P , w.h.p., the projected set has diam:

$$\text{diam}(PT) \leq \begin{cases} C \sqrt{\frac{m}{n}} \text{diam}(T), & \text{if } m \geq d(T) \\ C w_s(T), & \text{if } m \leq d(T). \end{cases}$$

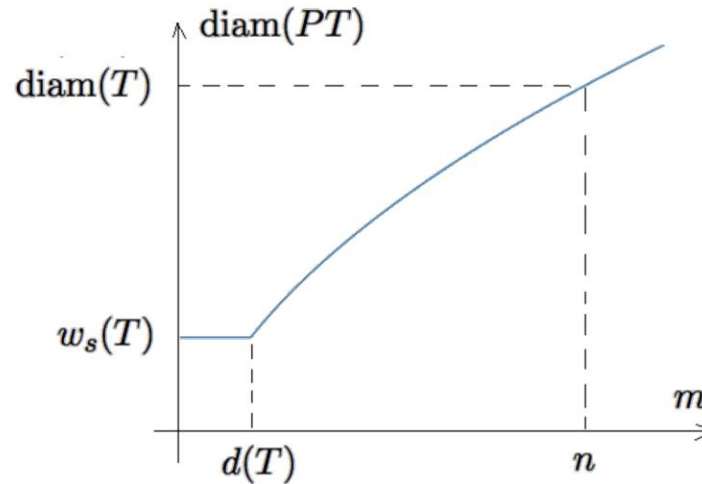


Figure 7.7 The diameter of a random m -dimensional projection of a set T as a function of m .

Take-away Messages

1. **Gaussian width and Gaussian complexity**
2. **Stable dimension and stable rank**
3. **Random projection (illustration of stable dimension)**

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Thanks!