

HDP

(25) Sparse Recovery

@ 滕佳焯

前情回顾

M^* bound: For the intersect of a set T and a random space, its diam can be very small w.h.p.

Escaping Theorem: with a large m , the intersection can be empty w.h.p.

Recovery Problem

Signal $x \in R^n$

(noisy measurement) $Y = Ax + w \in R^m$ ($m \ll n$)

Goal: recover x from (A, Y)

Different from regression, the goal is to recover x instead of estimating coefficient (but similarly).

Besides, we can usually choose A to meet some property (instead, the distribution of x is fixed in linear regression)

Note that when $m \ll n$, there can be many solutions (ill-posed).

Therefore, we need to add some prior information on signal x .

Noiseless Recovery Problem

Signal $x \in R^n$, (**noiseless** measurement) $Y = Ax - w \in R^m$ ($m \ll n$)

Prior information on x : $x \in T$

All the solutions forms a kernel space of A ($x \in x' + Ker(A)$)

Kernel space (with a shift)? M^* bound?

Theorem 10.2.1. *Suppose the rows A_i of A are independent, isotropic and sub-gaussian random vectors. Then any solution \hat{x} of the program (10.4) satisfies*

$$\mathbb{E} \|\hat{x} - x\|_2 \leq \frac{CK^2 w(T)}{\sqrt{m}},$$

where $K = \max_i \|A_i\|_{\psi_2}$.

| |
|---|
| $\text{find } x' : y = Ax', \quad x' \in T. \tag{10.4}$ |
|---|

Every point that satisfies $Ax=Y$ is close to the true signal x .

Noisy Recovery Problem

Signal $x \in R^n$, (**noisy** measurement) $Y = Ax + w \in R^m$ ($m \ll n$)

Prior information on x : $x \in T$

Exercise 10.2.4 (Noisy measurements). ☕☕ Extend the recovery result (Theorem 10.2.1) for the noisy model $y = Ax + w$ we considered in (10.1). Namely, show that

$$\mathbb{E} \|\hat{x} - x\|_2 \leq \frac{CK^2 w(T) + \|w\|_2}{\sqrt{m}}.$$

| | |
|--------------------------------------|--------|
| find $x' : y = Ax', \quad x' \in T.$ | (10.4) |
|--------------------------------------|--------|

Hint: although x' does not satisfy $Ax' = Ax$, it is close to x^0 (any point with $Ax^0 = Ax$, choose a proper one) since $A(x' - x^0) = w$

The prior T : Sparsity (noiseless)

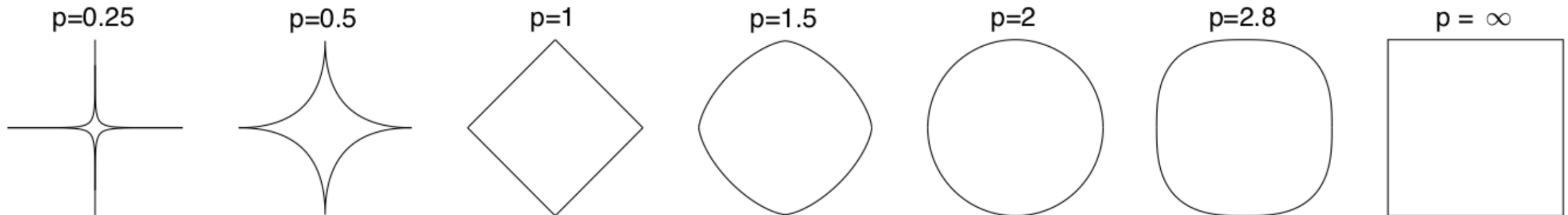
$$T = \{x \in \mathbb{R}^n : \|x\|_0 \leq s\}.$$

The problem is well posed when $m > 2\|x\|_0$

However, it is hard to calculate since T is non-convex.

Therefore, we want to relax T to its convex version.

$$\{x \in \mathbb{R}^n : \|x\|_0 \leq s, \|x\|_2 \leq 1\} \subset \sqrt{s}B_1^n.$$



The prior T: Sparsity (noiseless)

$$\{x \in \mathbb{R}^n : \|x\|_0 \leq s, \|x\|_2 \leq 1\} \subset \sqrt{s}B_1^n.$$

Corollary 10.3.4 (Sparse recovery: guarantees). *Assume the unknown s -sparse signal $x \in \mathbb{R}^n$ satisfies $\|x\|_2 \leq 1$. Then x can be approximately recovered from the random measurement vector $y = Ax$ by a solution \hat{x} of the program (10.6). The recovery error satisfies*

$$\mathbb{E} \|\hat{x} - x\|_2 \leq CK^2 \sqrt{\frac{s \log n}{m}}.$$

Note that the approximation is almost tight!
(the gaussian width is optimal up to a constant).

Low-rank Matrix Recovery

Signal $x \in \mathbb{R}^{d \times d}$, (**noiseless** measurement) $Y_i = A_i \cdot X \in \mathbb{R}^m$ ($m \ll n$)

Prior information on X : $\text{Rank}(X) \leq r$ [c.f. $\|x\|_0$ in vector]

Nuclear norm (c.f. $\|x\|_1$) $\|X\|_* := \|s(X)\|_1 = \sum_{i=1}^d s_i(X) = \text{tr}(\sqrt{X^\top X})$

$$\{X \in \mathbb{R}^{d \times d} : \text{rank}(X) \leq r, \|X\|_F \leq 1\} \subset \sqrt{r} B_*.$$

Exercise 10.4.5 (Low-rank matrix recovery: guarantees). ☕☕ Suppose the random matrices A_i are independent and have all independent, sub-gaussian entries.¹ Assume the unknown $d \times d$ matrix X with rank r satisfies $\|X\|_F \leq 1$. Show that X can be approximately recovered from the random measurements y_i by a solution \hat{X} of the program (10.11). Show that the recovery error satisfies

$$\mathbb{E} \|\hat{X} - X\|_F \leq CK^2 \sqrt{\frac{rd}{m}}.$$

Hint: use the Gaussian width of nuclear norm ball.

Exact sparse recovery (noiseless)

Theorem 10.5.1 (Exact sparse recovery). *Suppose the rows A_i of A are independent, isotropic and sub-gaussian random vectors, and let $K := \max_i \|A_i\|_{\psi_2}$. Then the following happens with probability at least $1 - 2 \exp(-cm/K^4)$.*

Assume an unknown signal $x \in \mathbb{R}^n$ is s -sparse and the number of measurements m satisfies

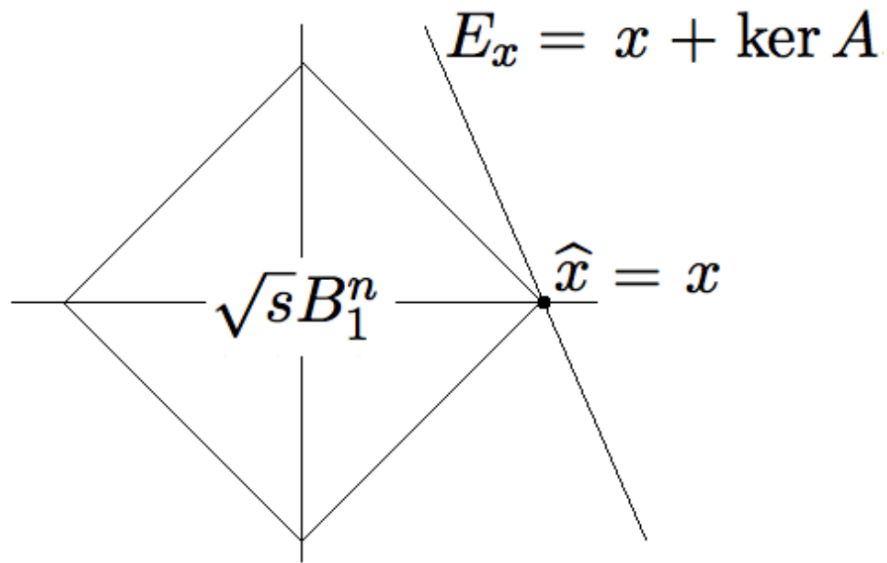
$$m \geq CK^4 s \log n.$$

Then a solution \hat{x} of the program (10.12) is exact, i.e.

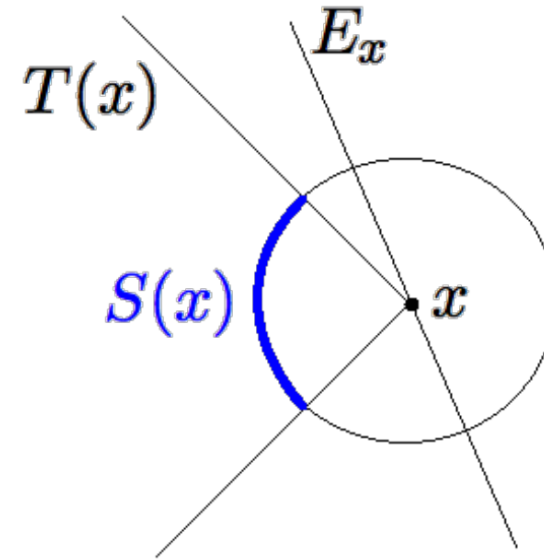
$$\hat{x} = x.$$

When m (#samples) is large, the signal can be recovered perfectly.

Exact sparse recovery (noiseless, intuition)



(a) Exact sparse recovery happens when the random subspace E_x is tangent to the ℓ_1 ball at the point x .



(b) The tangency occurs iff E_x is disjoint from the spherical part $S(x)$ of the tangent cone $T(x)$ of the ℓ_1 ball at point x .

Hint: prove that the residual $h = x - \hat{x}$ satisfying $\|h\|_1 \leq \sqrt{2s} \|h\|_2$

Exact sparse recovery (noiseless, general RIP condition)

Definition 10.5.8 (RIP). An $m \times n$ matrix A satisfies the *restricted isometry property* (RIP) with parameters α , β and s if the inequality

$$\alpha\|v\|_2 \leq \|Av\|_2 \leq \beta\|v\|_2$$

holds for all vectors $v \in \mathbb{R}^n$ such that² $\|v\|_0 \leq s$.

Intuition: the matrix A is approximate isometric, closely related to eigenvalues.

Theorem 10.5.10 (RIP implies exact recovery). Suppose an $m \times n$ matrix A satisfies RIP with some parameters α , β and $(1 + \lambda)s$, where $\lambda > (\beta/\alpha)^2$. Then every s -sparse vector $x \in \mathbb{R}^n$ can be recovered exactly by solving the program (10.12), i.e. the solution satisfies

$$\hat{x} = x.$$

Exact recovery (noiseless, general RIP condition)

Random matrix meets the RIP condition.

Theorem 10.5.11 (Random matrices satisfy RIP). *Consider an $m \times n$ matrix A whose rows A_i of A are independent, isotropic and sub-gaussian random vectors, and let $K := \max_i \|A_i\|_{\psi_2}$. Assume that*

$$m \geq CK^4 s \log(en/s).$$

Then, with probability at least $1 - 2 \exp(-cm/K^4)$, the random matrix A satisfies RIP with parameters $\alpha = 0.9\sqrt{m}$, $\beta = 1.1\sqrt{m}$ and s .

LASSO

$$\text{minimize } \|y - Ax'\|_2 \text{ s.t. } \|x'\|_1 \leq R.$$

Theorem 10.6.1 (Performance of Lasso). *Suppose the rows A_i of A are independent, isotropic and sub-gaussian random vectors, and let $K := \max_i \|A_i\|_{\psi_2}$. Then the following happens with probability at least $1 - 2 \exp(-s \log n)$.*

Assume an unknown signal $x \in \mathbb{R}^n$ is s -sparse and the number of measurements m satisfies

$$m \geq CK^4 s \log n. \tag{10.23}$$

Then a solution \hat{x} of the program (10.22) with $R := \|x\|_1$ is accurate, namely

$$\|\hat{x} - x\|_2 \leq C\sigma \sqrt{\frac{s \log n}{m}},$$

where $\sigma = \|w\|_2 / \sqrt{m}$.

Remark 10.6.3 (Exact recovery). In the noiseless model $y = Ax$ we have $w = 0$ and thus Lasso recovers x exactly, i.e.

$$\hat{x} = x.$$

Take-away Messages (noiseless setting)

Vector recovery problem:

- $\ell_0 + \ell_2 \approx \ell_1$
- Error bound $\sqrt{s \log n / m}$ [M* bound]
- Exact recovery $m \sim s \log n$ [Escaping Theorem]

Matrix recovery problem:

- $\|\cdot\|_0 + \|\cdot\|_F \approx \|\cdot\|_*$
- RIP condition (random matrix)
- Error bound $\sqrt{rd/m}$

LASSO (noisy):

- Error bound $\sigma \sqrt{s \log n / m}$

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Thanks!